





CSP-based inference of function block finite-state models from execution traces

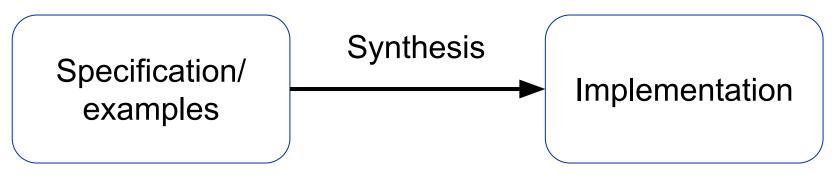


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25 July 2017

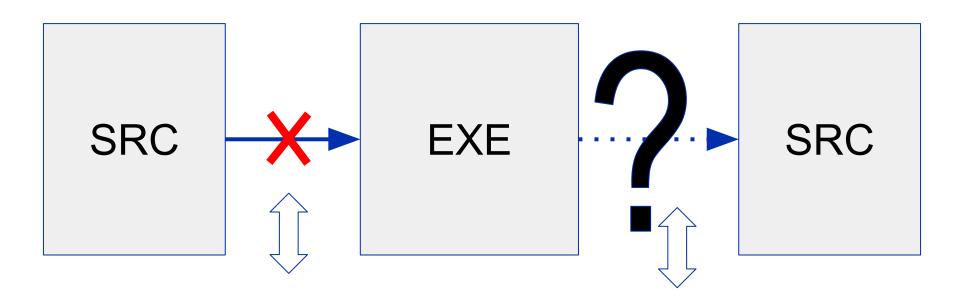
Program synthesis

- Derive implementation from examples/specification
 - From seminal work [A. Church, 1963]



- Motivation
 - Fundamental in computer science
 - Automation of software engineering
 - Reverse engineering

Reverse engineering of software



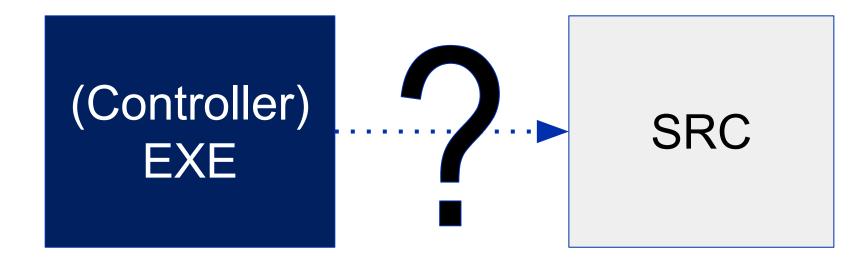
- Rights limitations
- Changing standards

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- Understanding
- Optimization
- Verification

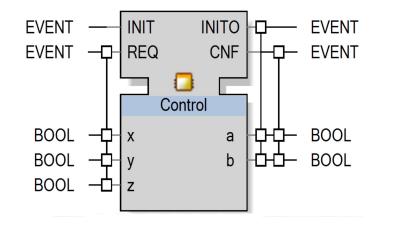


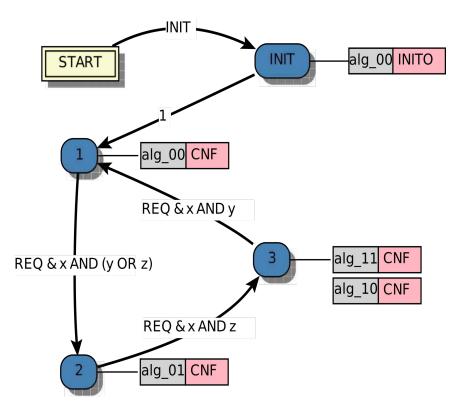
Black-box approach



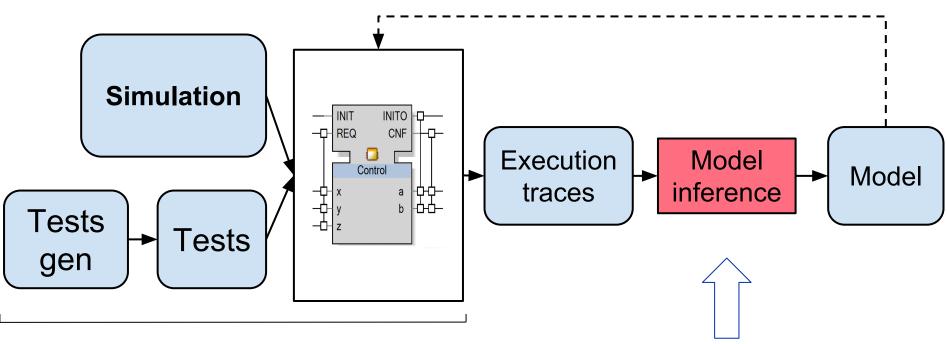
What's in the box?

Target language: IEC 61499 function blocks



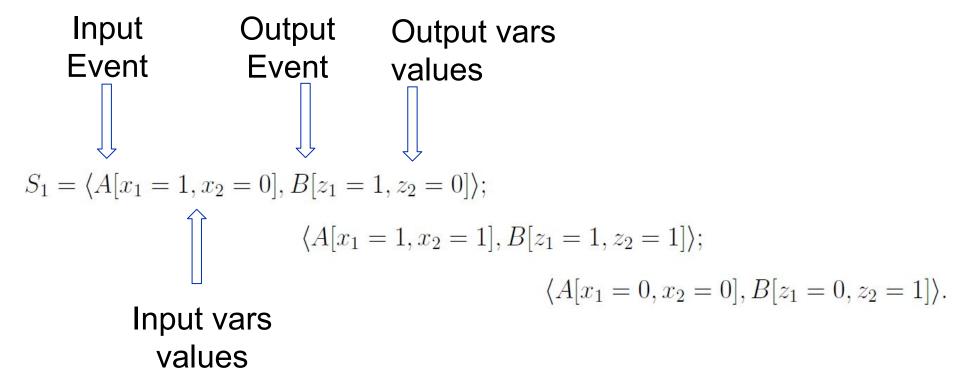


Test-based reverse engineering



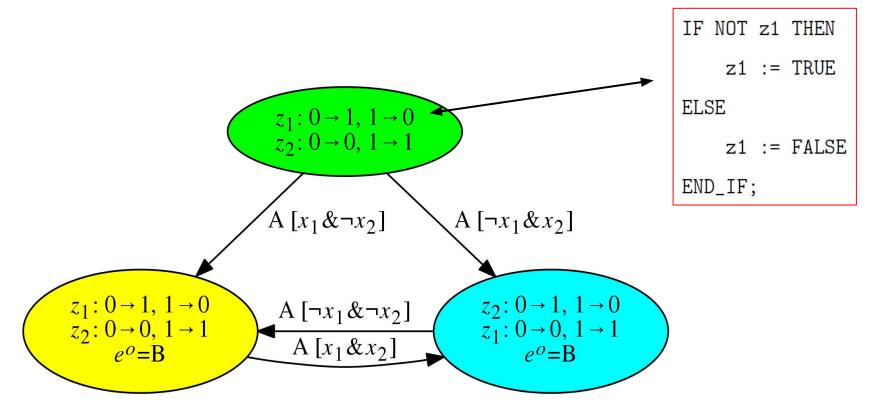
Preparation

Scenarios



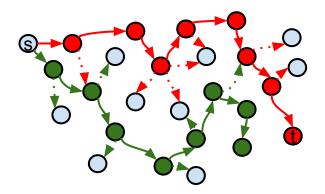
Basic function block model

Boolean input/output vars



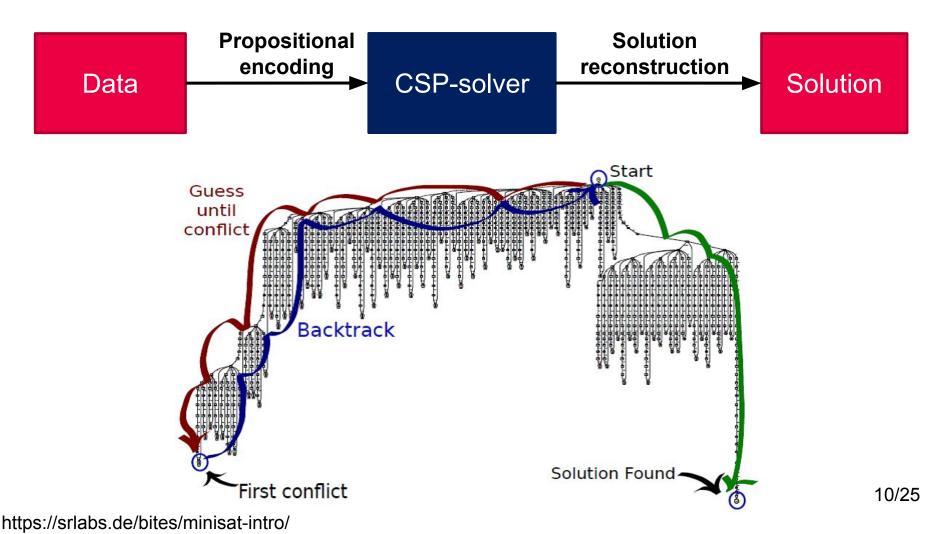
Previous/proposed approaches

- 1. Metaheuristic: [Chivilikhin et al / INDIN'15]
 - Slow
 - Approximate



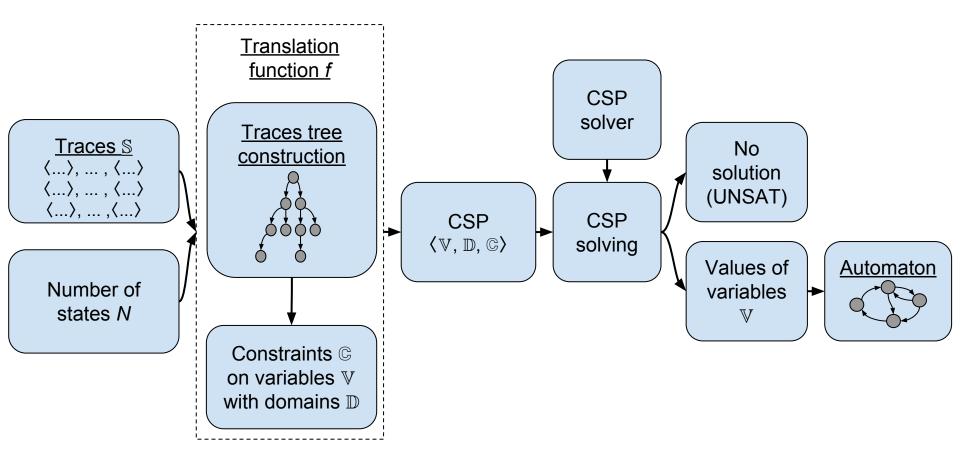
- 2. We propose CSP-translation approach
 - Could be faster in practice
 - Exact

Proposed approach: translation to Constraint Satisfaction Problem





Proposed approach scheme



Traces tree

$$S_{1} = \langle A[x_{1} = 1, x_{2} = 0], B[z_{1} = 1, z_{2} = 0] \rangle;$$

$$\langle A[x_{1} = 1, x_{2} = 1], B[z_{1} = 1, z_{2} = 1] \rangle;$$

$$S_{2} = \langle A[x_{1} = 0, x_{2} = 1], B[z_{1} = 0, z_{2} = 1] \rangle;$$

$$\langle A[x_{1} = 0, x_{2} = 0], B[z_{1} = 1, z_{2} = 1] \rangle;$$

$$\langle A[x_{1} = 1, x_{2} = 1], B[z_{1} = 1, z_{2} = 0] \rangle.$$

$$(A[x_{1} = 1, x_{2} = 1], B[z_{1} = 1, z_{2} = 0] \rangle.$$

$$e^{in} = A$$

$$x_{1} = 1$$

$$x_{2} = 0$$

$$e^{in} = A$$

$$x_{1} = 1$$

$$x_{2} = 0$$

$$x_{1} = 1$$

$$x_{2} = 0$$

$$x_{2} = 0$$

$$x_{1} = 1$$

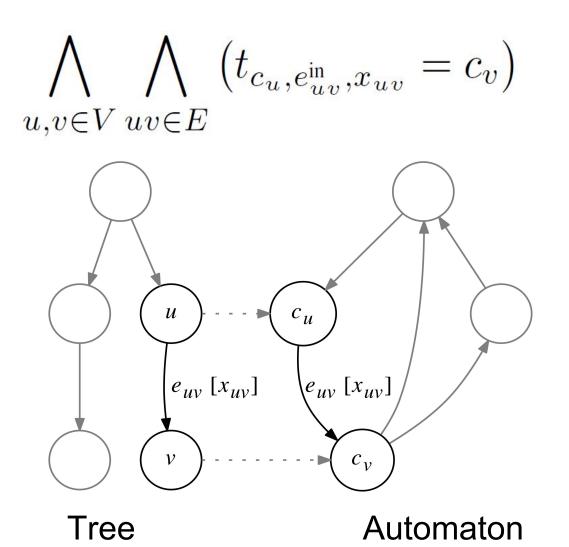
$$x_{2} = 1$$

$$x_{3} = 1$$

Variables

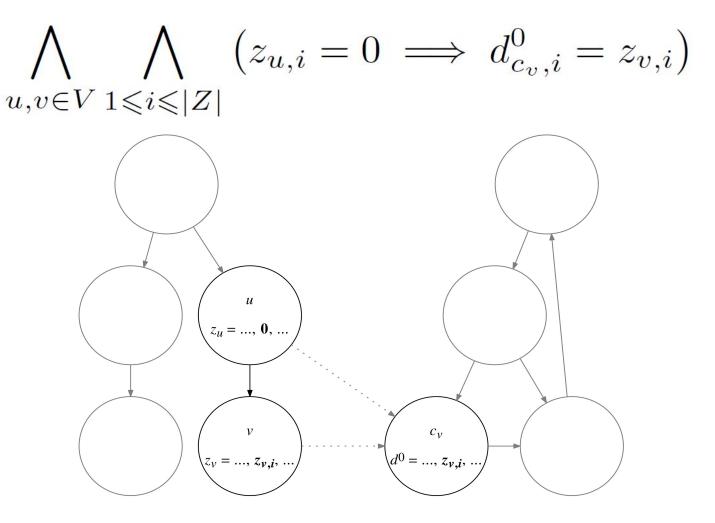
- $c_v \in [1..N]$ for $v \in V$ color of traces tree vertex v;
- $t_{n,e,x} \in [1..N]$ for $n \in [1..N]$, $e \in [1..|E^{I}|]$, $x \in [1..|\hat{X}|]$ target state of the transition from state *n* labeled with input event *e* and guard condition *x*;
- $o_n \in [1..|E^{O}| + 1]$ for $n \in [1..N]$ index of the output event in state $n (|E^{O}| + 1 \text{ corresponds to } \varepsilon)$;
- d⁰_{n,i} ∈ {0,1} for n ∈ [1..N], i ∈ [1..|Z|] value of the *i*-th output variable in state n if its previous value equals zero;
- $d_{n,i}^1 \in \{0,1\}$ $n \in [1..N], i \in [1..|Z|]$ value of the *i*-th output variable in state n if its previous value equals one.

Main constraints (1)



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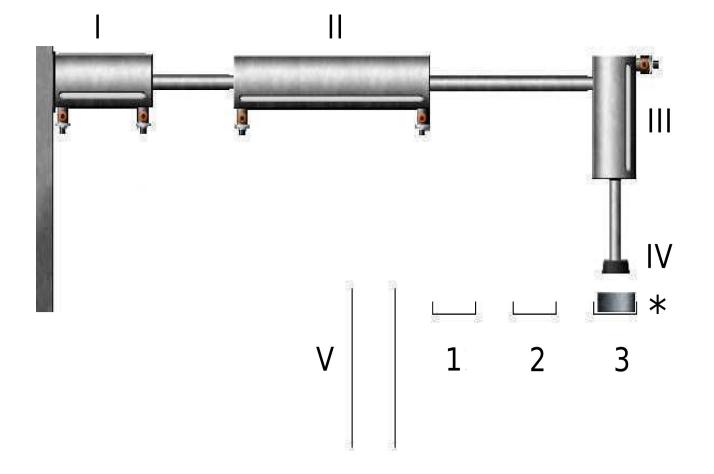
Main constraints (2)



Tree

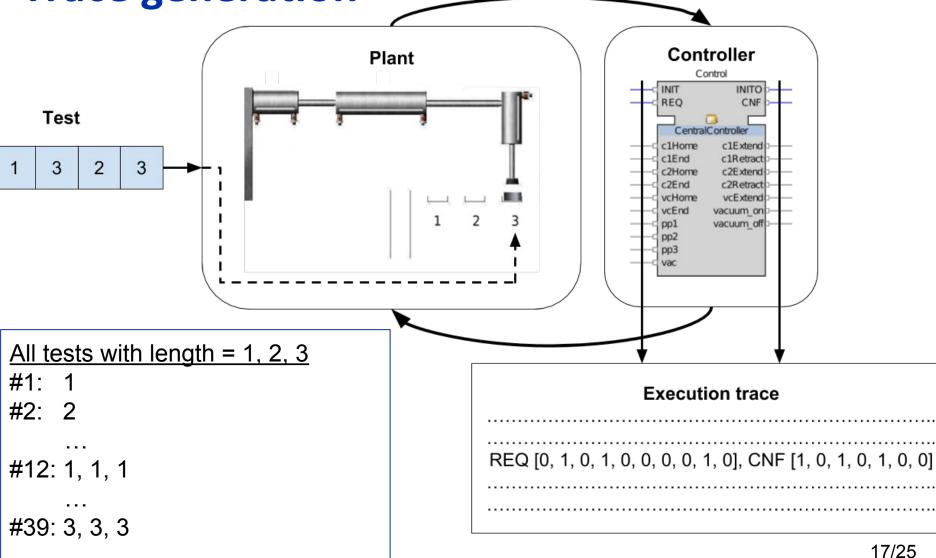
Automaton

Case study: pick-and-place manipulator





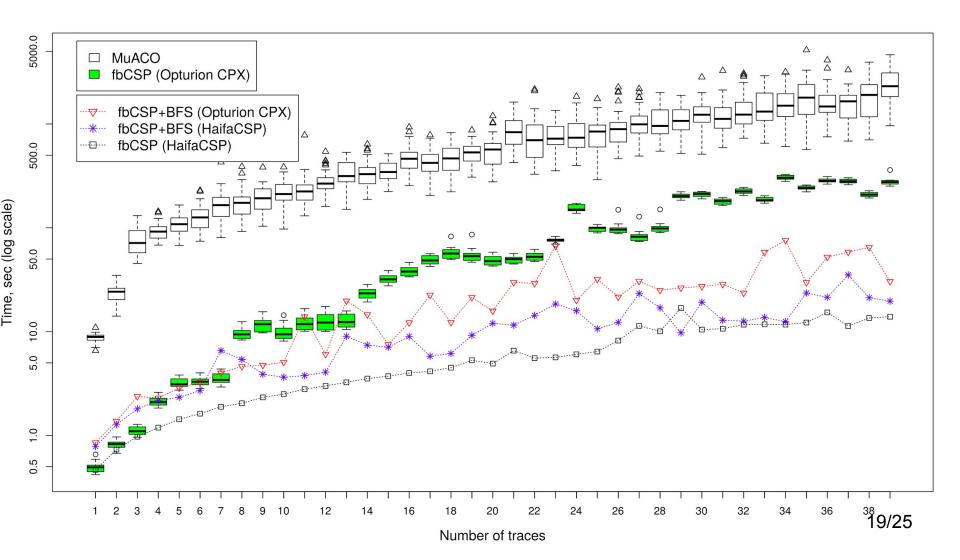
Trace generation



Experimental setup

- Methods
 - MuACO
 - metaheuristic [Chivilikhin et al (2015)]
 - ✓ fbCSP
 - Proposed approach
 - ✓ fbCSP+BFS
 - fbCSP + BFS-based symm breaking
- State-of-the-art CSP-solvers
 - ✓ Opturion CPX (Minizinc Challenge 2015 winner)
 - HaifaCSP (Minizinc Challenge 2016 winner)

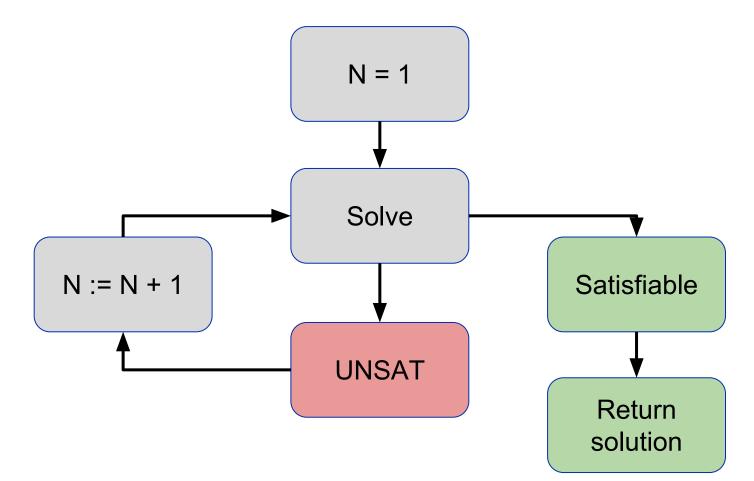
Fixed number of states: N = 10



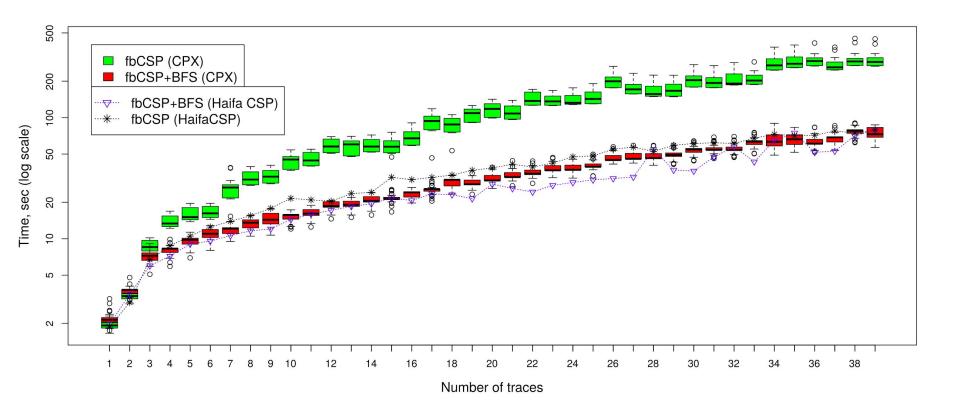
Minimal model generation

- Most general pattern in the given data
- •Occam's razor (law of parsimony):
 - "Among competing hypotheses, the one with the fewest assumptions should be selected"

Minimal model generation

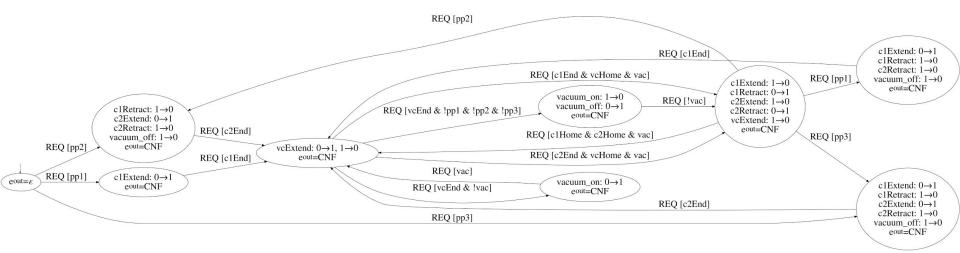


Minimal model generation: results





Generated model example



Conclusion and Future work

- Proposed fast exact algorithm for inferring minimal-sized models of basic FBs for logic control
- Available: https://github.com/chivdan/cspgen
- Future/ongoing work
 - Integer/real variables
 - Timers
 - Composite FBs
 - CEGAR for LTL/CTL based inference



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Thank you for your attention!

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